# Is there a group with order less than 90 that can realize $6 \times 6$ matrix multiplication better than Strassen's algorithm?

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# 1 Problem statement

**Problem Restatement:** The problem is equal to "Is there a group with order less than 90 that can realize < 6, 6, 6 > TPP property and have multiplication rank less than 161?[1]".

### 2 Method to solve it

Since the search space is too large, my main thinking is to reduce the search space by lots of necessary conditions.

#### 2.1 necessary conditions and how they reduce the search space

For a finite group G, let T(G) be the number of irreducible complex characters of G and b(G) the largest degree of an irreducible character of G.

**Theorem 2.1** ([2). , Theorem 6 and Remark 2] Let G be a group. (1) If b(G) = 1, then R(G) = |G|. (2) If b(G) = 2, then R(G) = 2|G| - T(G). (3) If  $b(G) \ge 3$ , then  $R(G) \ge 2|G| + b(G) - T(G) - 1$ .

**Definition 1** (Triple Product Property). We say that the nonempty subsets S, T and U of a group G satisfy the Triple Product Property(*TPP*) if for  $s \in Q(S)$ ,  $t \in Q(T)$  and  $u \in Q(U)$ , stu=1 holds if and only if s=t=u=1. If this holds, we say the the group G realizes  $\langle |S|, |T|, |U| \rangle$  via S, T, U.

**Definition 2.** Let  $\beta(G)$  be the maximum of  $n^*m^*p$ , where G realizes  $\langle n, m, p \rangle$ .

**Theorem 2.2.** For an abelian group,  $\omega = 3$ .

*Proof.* If G is abelian and non-trivial  $(|G| \neq 1)$ , then b(G)=1 and from Theorem 1 we have: R(G) = |G|. And from page 3 of [3], we know that  $R(n, m, p) \leq R(G)$ . So  $\beta(G) \leq R(G) = |G|$ . From Theorem 1.7. in [3], we have  $\omega \leq 3$ .

**Remark 2.1.** Then Since we are looking for non-trivial solutions for  $\omega$ , now we only need to consider non-abelian groups.

Lemma 2.1. (abelian judgements)

(1) If |G| is a prime then G is abelian.

(2) If |G| = pq, where p and q are primes, p;q, if  $q \not\equiv 1 \mod p$ , then G is abelian.

(3) If  $|G| = pq^2$ , p and q are two distinct primes and p doesn't divide -Aut(G), then G is abelian.

(4) If |G| = pqr, p, q and r are three distinct primes and q;r,  $r \not\equiv 1 \pmod{p}$ , qr; p,  $p \not\equiv 1 \pmod{p}$ ,  $p \not\equiv 1 \pmod{p}$ , then G is abelian.

**Theorem 2.3.** ([3]) If G is non-abelian, then  $T(G) \leq (5/8)|G|$ . Equality implies that |G : Z(G)| = 4.

**Remark 2.2.** Then if we combine Theorem 2 and Theorem 3, we have:  $R(G) \ge 2|G| - T(G) \ge (11/8)|G|$ . Since we want  $R(G) \le 161$ , then we have: (11/8)|G| < 161 $|G| \le 117$ .

**Definition 3.** (C1 candidates, similar but not quite the same to [3], Definition 3.2) A group G that realizes < 6, 6, 6 > and satisfies  $\underline{R}[G] < 161$  will be called C1 candidate.

**Theorem 2.4.** ([4], Observation 3.1) Let (s,t,u) be the parameters of a TPP triple in G. Then  $s(t+u-1) \leq |G|, t(s+u-1) \leq |G|$  and  $u(s+t-1) \leq |G|$ .

**Proposition 1.** If G is a C1 candidate, then  $66 \le |G| \le 117$ .

*Proof.* From the theorem above, we have  $|G| \ge 6 * (6 * 6 - 1) = 66$ . Consider Theorem 1 and Theorem 4 above, we have  $R(G) \ge 2|G| - T(G) \ge (11/8)|G|$ . And  $\underline{R}[G] < 161$ , then we have  $|G| \le 117$ .

**Remark 2.3** (**GAP experiment**). After the "abelian judgement (1), (2) and (4)" stated above, if G is a C1 candidate, then  $|G| \in \{66, 68, 70, 72, 74, 75, 76, 78, 80, 81, 82, 84, 86, 88, 90, 92, 93, 94, 96, 98, 99, 100, 102, 104, 105, 106, 108, 110, 111, 112, 114, 116, 117\}.$ 

**Definition 4.** ([5], Definition 3.4) Let G be a group with a TPP triple (S, T, U), and suppose H is a subgroup of index 2 in G. We define  $S_0 = S \cap H$ ,  $T_0 = T \cap H$ ,  $U_0 = U \cap H$ ,  $S_1 = S \setminus H$ ,  $T_1 = T \setminus H$  and  $U_1 = U \setminus H$ .

**Lemma 2.2.** Suppose G realizes < 6, 6, 6 >. If G has a subgroup H of index 2, then H realizes < 3, 3, 3 >.

*Proof.* Suppose G realizes  $\langle 6, 6, 6 \rangle$  with the TPP triple (S,T,U). If  $|S_0| \langle S_1|$ , then for any  $a \in S_1$ , replace S by  $Sa^{-1}$ . This will have the effect of interchanging  $S_0$  and  $S_1$ . Hence we may assume that  $|S_0| \geq |S_1|, |T_0| \geq |T_1|$  and  $|U_0| \geq |U_1|$ . Now  $(S_0, T_0, U_0)$  is a TPP triple of H, and since each of  $S_0, T_0$  and  $U_0$  has at least 3 elements, then H realizes  $\langle 3, 3, 3 \rangle$ .

**Lemma 2.3.** ([3],Lemma 3.6) Suppose G has a TPP triple (S,T,U). Let H be an abelian subgroup of index 2 in G. Then the following hold.  $a)|S_0^{-1}T_0U_0| = |S_0||T_0||U_0|;$  $b)|S_1^{-1}T_1U_0| \ge |S_1||T_1|;$   $\begin{array}{l} c)|S_1^{-1}U_1| = |S_1||U_1|;\\ d)S_0^{-1}T_0U_0 \cap S_1^{-1}T_1U_0 = \emptyset;\\ e)S_0^{-1}T_0U_0 \cap S_1^{-1}U_1T_0 = \emptyset;\\ f)S_1^{-1}T_1U_0 \cap S_1^{-1}U_1T_0 = \emptyset. \end{array}$ 

**Lemma 2.4.** If G realizes < 6, 6, 6 > and |G| < 90, then G has no abelian subgroups of index 2.

*Proof.* Suppose G has an abelian subgroup H of index 2 and realizes  $\langle 6, 6, 6 \rangle$  via the TPP triple (S,T,U). Define  $S_0, T_0, U_0, S_1, T_1, U_1$  as before. Then, as proved above, we may assume  $|S_0| \geq 3$ ,  $|T_0| \geq 3$  and  $|U_0| \geq 3$ . Without loss of generality we may assume that  $|S_0| \geq |T_0|$  and  $|S_0| \geq |U_0|$ .Now since  $|G| \leq 95$ , then  $|H| \leq 47$ . From the last lemma, we have  $47 \geq |H| \geq |S_0^{-1}T_0U_0 \cup S_1^{-1}U_1T_0 \cup S_1^{-1}T_1U_0|$ 

$$\begin{split} &= |S_0||T_0||U_0| + |S_1^{-1}U_1T_0| + |S_1^{-1}T_1U_0| \quad (3) \\ &\geq |S_0||T_0||U_0| + |S_1||U_1| + |S_1||T_1|. \quad (4) \\ &\text{If } |U_0| \geq 4, \text{ then } |H| \geq 64, \text{contradiction.So } |U_0| = 3: \\ &(a)|U_0| = 3 = |T_0| = |S_0|, \text{then from } (4) \text{we have: } |H| \geq 45, \text{contradiction.} \end{split}$$

(b) $|U_0| = 3 = |T_0|, |S_0| = 4$ , then from (4) we have:  $|H| \ge 48$ , contradiction.

 $(c)|U_0| = 3, |T_0| = |S_0| = 4$ , then from (4) we have  $|H| \ge 58$ , contradiction.

(d) $|U_0| = 3 = |T_0|, |S_0| = 5$ , then from (4) we have  $|H| \ge 51$ , contradiction.

(e) $|U_0| = 3, |T_0| = 4, |S_0| = 5$ , then from (4) we have  $|H| \ge 65$ , contradiction.

 $(f)|U_0| = 3, |T_0| = |S_0| = 5$ , then from (4) we have  $|H| \ge 79$ , contradiction.

 $(g)|U_0| = 3 = |T_0|, |S_0| = 6$ , then from (4) we have  $|H| \ge 54$ , contradiction.

(h) $|U_0| = 3$ ,  $|T_0| = 4$ ,  $|S_0| = 6$ , then from (4) we have  $|H| \ge 72$ , contradiction.

(i) $|U_0| = 3, |T_0| = 5, |S_0| = 6$ , then from (4) we have  $|H| \ge 90$ , contradiction.

 $(j)|U_0| = 3, |T_0| = |S_0| = 6$ , then from (4) we have  $|H| \ge 108$ , contradiction.

**Remark 2.4.** Up to now, among groups of order  $\leq 89$ , we have these 56 left for C1 candidates: (68,3),(72,3),(72,15),(72,16),(72,19),(72,20),(72,21),(72,22),(72,23),(72,24),(72,25),(72,39),(72,40),(72,41),(72,42),(72,43),(72,44),(72,45),(72,46),(72,47),(75,2),(78,1),(78,2),(80,3),(80,15),(80,18),(80,28),(80,29),(80,30),(80,31),(80,32),(80,33),(80,34),(80,39),(80,40),(80,41),(80,42),(80,49),(80,50),(81,3),(81,4),(81,6),(81,7),(81,8),(81,9),(81,12),(81,13),(81,14),(84,1),(84,2),(84,7),(84,8),(84,9),(84,10),(84,11).

**Theorem 2.5.** ([5], Theorem 1.8) Suppose G realizes  $\langle n, m, p \rangle$  and the character degrees of G are  $\{d_i\}$ . Then  $(nmp)^{\omega/3} \leq \sum d_i^{\omega}$ .

**Proposition 2.** The above theorem yields a nontrivial bound on  $\omega$  if and only if  $(nmp)^{\omega/3} \geq \sum d_i^{\omega}$ .

**Remark 2.5.** Since we are using the inequality stated in the proposition, in order to make  $\omega$  nontrivial, we need to search for groups has  $\sum d_i^{\omega} < 216$ .

**Remark 2.6.** (GAP experiment) So up to now, with the help of GAP experiment, we have these 18 groups(listed in their GAP ID) left as C1 candidates when its order < 90:(72,3),(72,16),(72,20),(72,21),(72,22),(72,23),(72,24),(72,25),(72,42),(72,46),(72,47),(81,3),(81,4),(81,6),(81,12),(81,13),(81,14),(84,10).

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