

Formally Specified Computer Algebra Software DK10

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Outline

- 1 Introduction
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Introduction

- Project goals
 - Formal specification of programs written in untyped computer algebra languages
 - Especially to find errors/inconsistencies
 - for example violation of method preconditions
- Computer algebra software at RISC as examples
 - DK11: rational parametric algebraic curves ([Maple](#))
 - DK6: computer algebra tools for special functions in numerical analysis ([Mathematica](#))
 - DK1: automated theorem proving ([Mathematica](#))

Past Activities (October 2009 to February 2010)

- Course work
 - Computer Algebra
 - Automated Theorem Proving
 - Formal Methods in Software Development
 - Thinking, Speaking, Writing
 - Formal Methods Seminar
 - Programming Project
- Literature study
 - Type systems
 - Polymorphism
 - Abstract data types
 - Denotational semantics
 - Functional programming languages
 - Pattern matching
 - Type checking and inference

Software Study - Computer Algebra

- Bivariate difference-differential dimension polynomials and their computation
- Relative Gröbner bases computation (using M. Zhou and F. Winkler's algorithm)
- Maple implementation of the algorithms
- Software
 - Maple package *DifferenceDifferential*
 - Christian Dönch
- Literature reference
 - Christian Dönch. *Bivariate difference-differential dimension polynomials and their computation in Maple*. Technical report no. 09-19 in RISC Report Series, University of Linz, Austria, 2009.

Software Study - Computer Algebra

```
ddsub := proc(c,b)  
    local f, g, i, m, n, a1;  
    f := c; g := b;  
for i to nops(g) do  
    g [i][1] := -g [i][1];  
    f := [op(f),g [i];  
end do;  
for m from nops(f) by -1 to 1 do  
    for n from m-1 by -1 to 1 do  
        if f [m][2] = f [n][2] and f [m][3] = f [n][3] and f [m][4] = f [n][4] then  
            a1 := f [m][1] + f [n][1]; f [n] := subsop(1=a1,f [n]);  
            f := subsop(m=NULL,f); n := m;  
        end if;  
    end do;  
end do;  
....  
return f;  
end proc;
```

Software Study - Computer Algebra

- Potential considerations
 - Limited types used i.e. integer and list
 - Not much use of Maple libraries - mostly standalone
 - No destructive update of data structures
 - Imperative style of development

Procedural/functional Maple package

Software Study - Algorithmic Combinatorics

- Advanced applications of holonomic systems approach
- Computations in Ore algebras
- Non-commutative Gröbner bases
- Solving linear system of differential equations
- Software
 - Symbolic summation and integration for holonomic functions
 - Mathematica package - *HolonomicFunctions*
 - Christoph Koutschan
- Literature reference
 - Christoph Koutschan. *HolonomicFunctions (User's Guide)*. Technical report no. 10-01 in RISC Report Series, JKU, Austria, January 2010.
 - Christoph Koutschan. *Advanced Applications of the Holonomic Systems Approach*. RISC-Linz, JKU. PhD Thesis, September 2009.

Software Study - Algorithmic Combinatorics

```
OrePlus [p1:OrePolynomial[data1_List, algebra_OreAlgebraObject, order_],
        p2:OrePolynomial[data2_List, algebra_OreAlgebraObject, order_]] :=
Module[{i1, i2, l1, l2, c, c1, c2, m1, m2, sum, coeffPlus},
  l1 = Length[data1];
  If[l1 === 0, Return[p2]];
  l2 = Length[data2];
  If[l2 === 0, Return[p1]];
  coeffPlus = algebra[[3]];
  i1 = 1; i2 = 1;
  sum = {};
  While[i1 <= l1 && i2 <= l2,
    {c1, m1} = data1[[i1]];
    {c2, m2} = data2[[i2]];
    If [m1 === m2, c = coeffPlus[c1, c2];
    If[Not[MatchQ[c, 0|0.]], AppendTo[sum, {c, m1}]];
    i1++; i2++; ,
  If[OreOrderedQ[m1, m2, order], AppendTo[sum, {c1, m1}]; i1++;
  .... ];];
```

Software Study - Algorithmic Combinatorics

- Potential considerations
 - Based on pattern matching
 - Imperative style of programming
 - Use of abstract data types
 - Use of customized Mathematica functionality
 - Not much use of Mathematica libraries

Procedural/functional Mathematica program with abstract data types

Software Study - Automated Theorem Proving

- Theorema set theory prover (STP)
- Automated prover for theorems
- Works with Prove-Compute-Solve (PCS) strategy
- Integrated with Theorema infrastructure (not standalone)
- Software
 - Mathematica package *SetTheoryProver*
 - Wolfgang Windsteiger
- Literature reference
 - W. Windsteiger. *An Automated Prover for Zermelo-Fraenkel Set Theory in Theorema*. JSC 41(3-4), pp. 435-470, 2006, Elsevier, ISSN 0747-7171.
 - W. Windsteiger. *A Set Theory Prover in Theorema: Implementation and Practical Applications*. RISC. PhD Thesis, May 2001.

Software Study - Automated Theorem Proving

```

STP[•lf[l_, T_ ∈ Intersection[A_, B_], i_], a_ • asml, af_] :=
  Module[
    {goalList=MapIndexed[•lf[NewLabel[∕MembershipAlternatives, l, #2], T_ ∈ #1, i] &,
    {A, B}], proofSits},
    proofSits=Map[Psit[#, a, af] &, goalList];
    ProofStep[Prinfo[∕MembershipFiniteIntersection, l, goalList],
    Sequence @@ proofSits]]
  
```

Software Study - Automated Theorem Proving

- Potential considerations
 - Pattern matching rules used
 - Are the rules overlapping?
 - Are the rules exhaustive?
 - Implicit type definitions
 - Declarative style of programming

Functional Mathematica program, essentially based on pattern matching

Current Activities

- Definition of simplified/typed versions of Mathematica and Maple (say *MiniMma* and *MiniMaple*)
 - Syntax
 - Type system
 - Semantics and soundness of typing
- Implementation of a type checker prototype
 - Static typing as a prerequisite to logic specification
- Experiments with software fragments available at RISC
- Next - Formal Specification language

Thanks for your attention!